Computing Precise Control Interface Specifications

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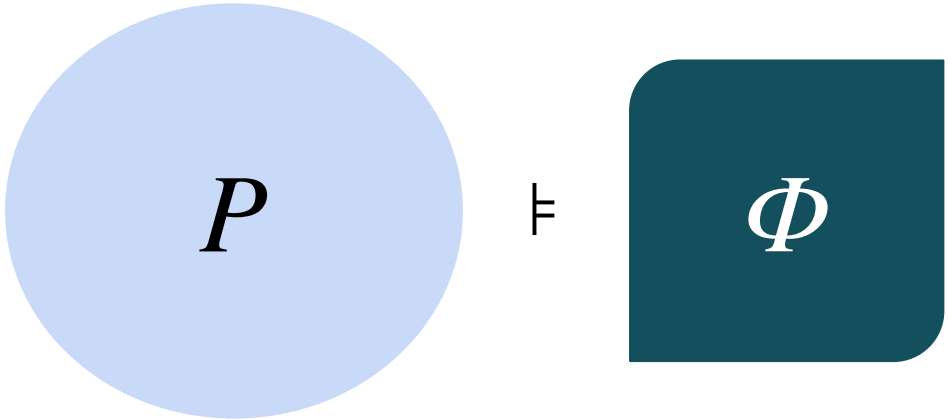
The Unknown

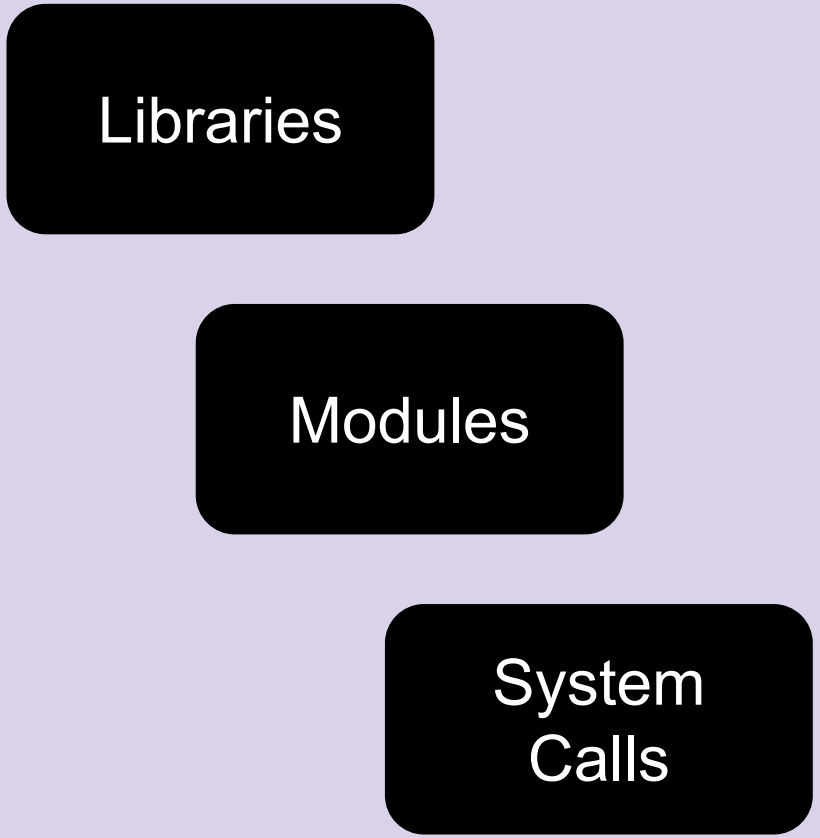
How do we verify programs with unknown code?

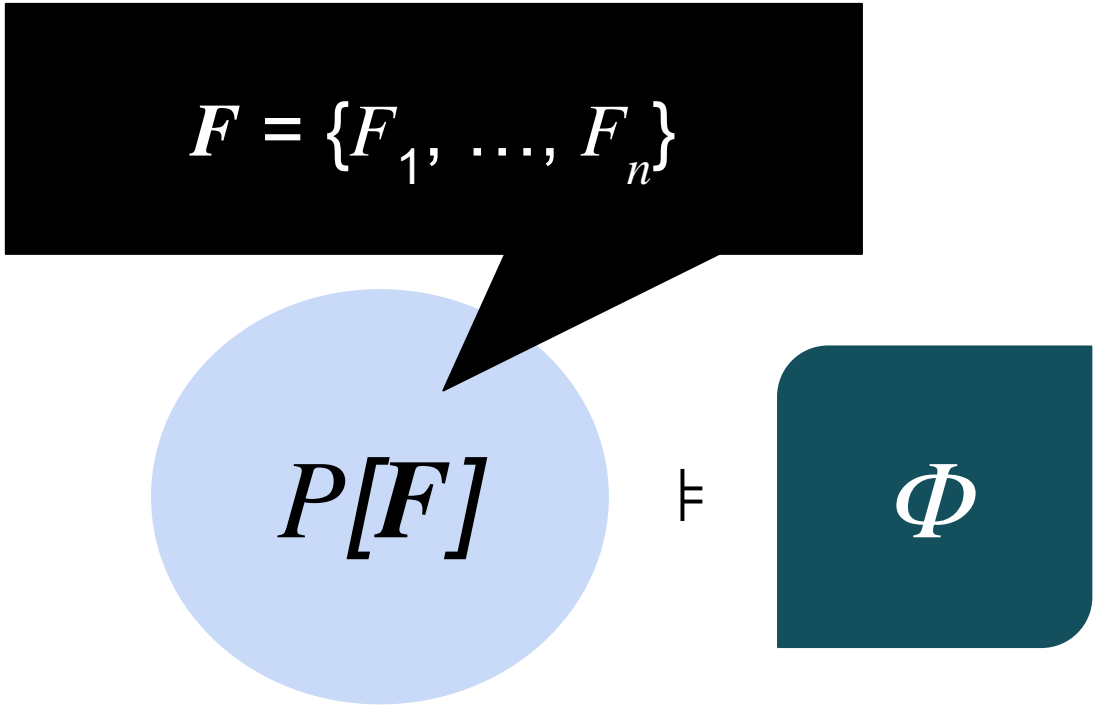
A Verification Myth

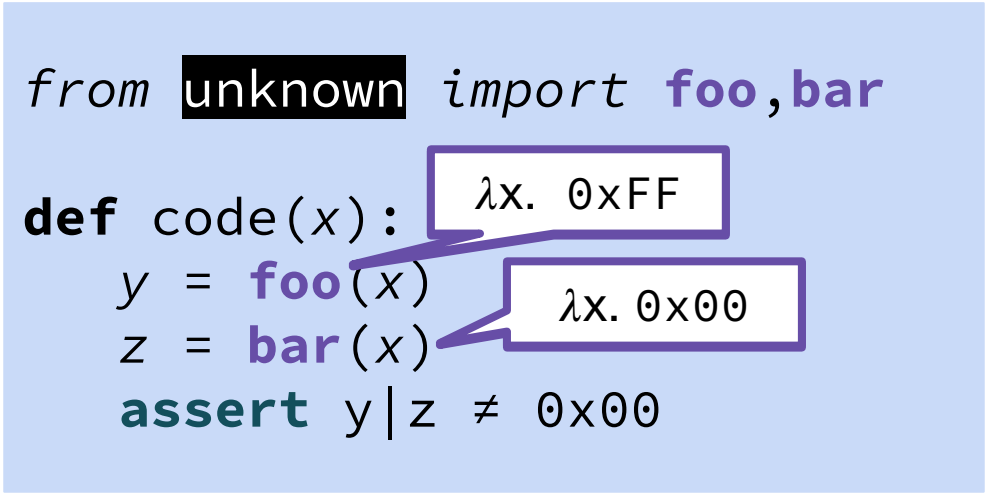


Program Spec

Source code is incomplete!




$F = \{F_1, \ldots, F_n\}$
$P[F]$
$\models$
Φ
Program
Spec



Which implementations?

```
import unknown as foo,bar

def code(x):
    y = foo(x)
    z = bar(x)
    assert y|z ≠ 0x00
```

Which implementations satisfy the spec?

Independent Specs
[POPL ‘16, SIGCOMM ‘20]

φ(**foo**) $\wedge$ ψ(**bar**)

Necessary [SIGCOMM ‘20]

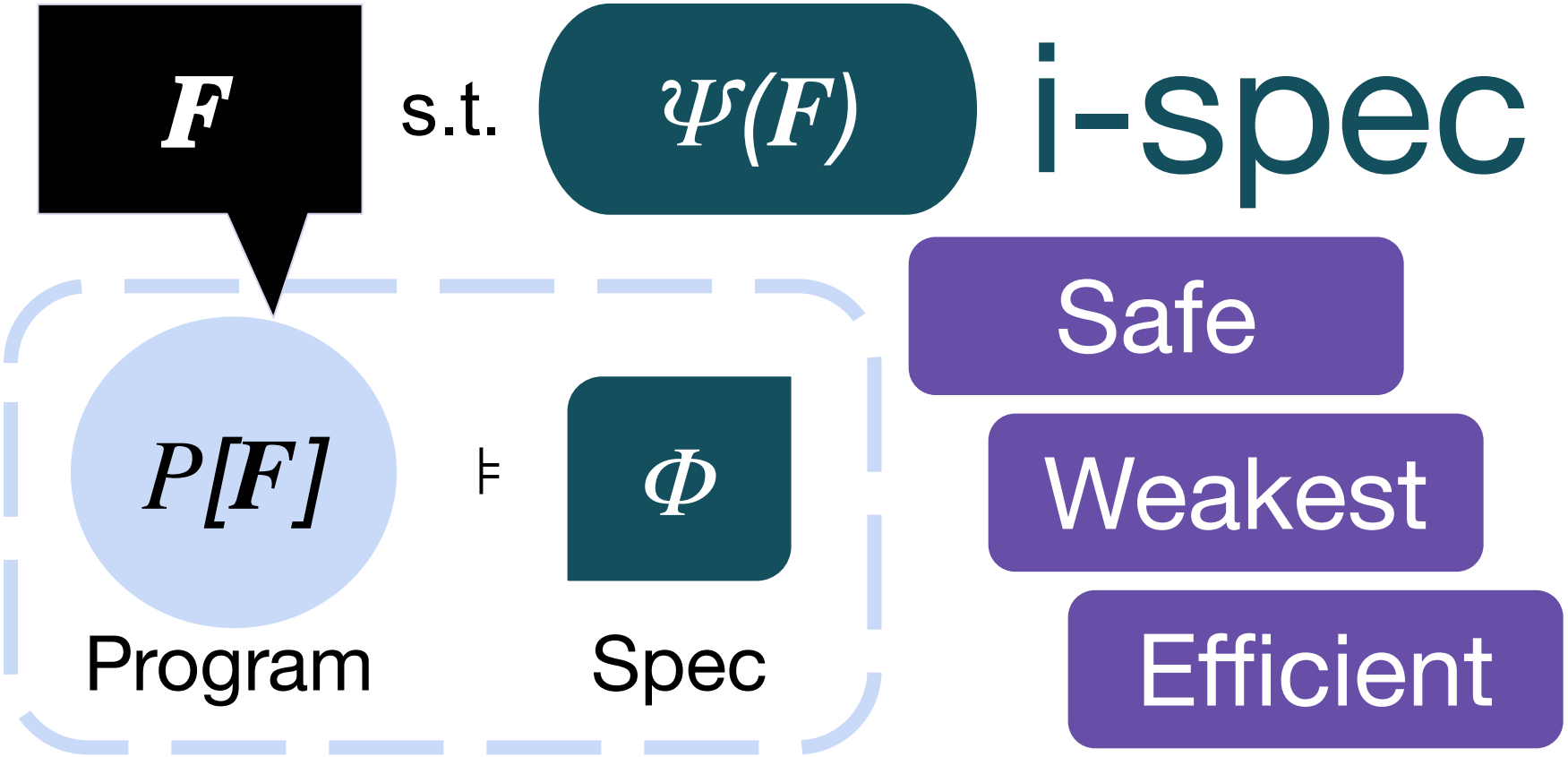
```
true
```

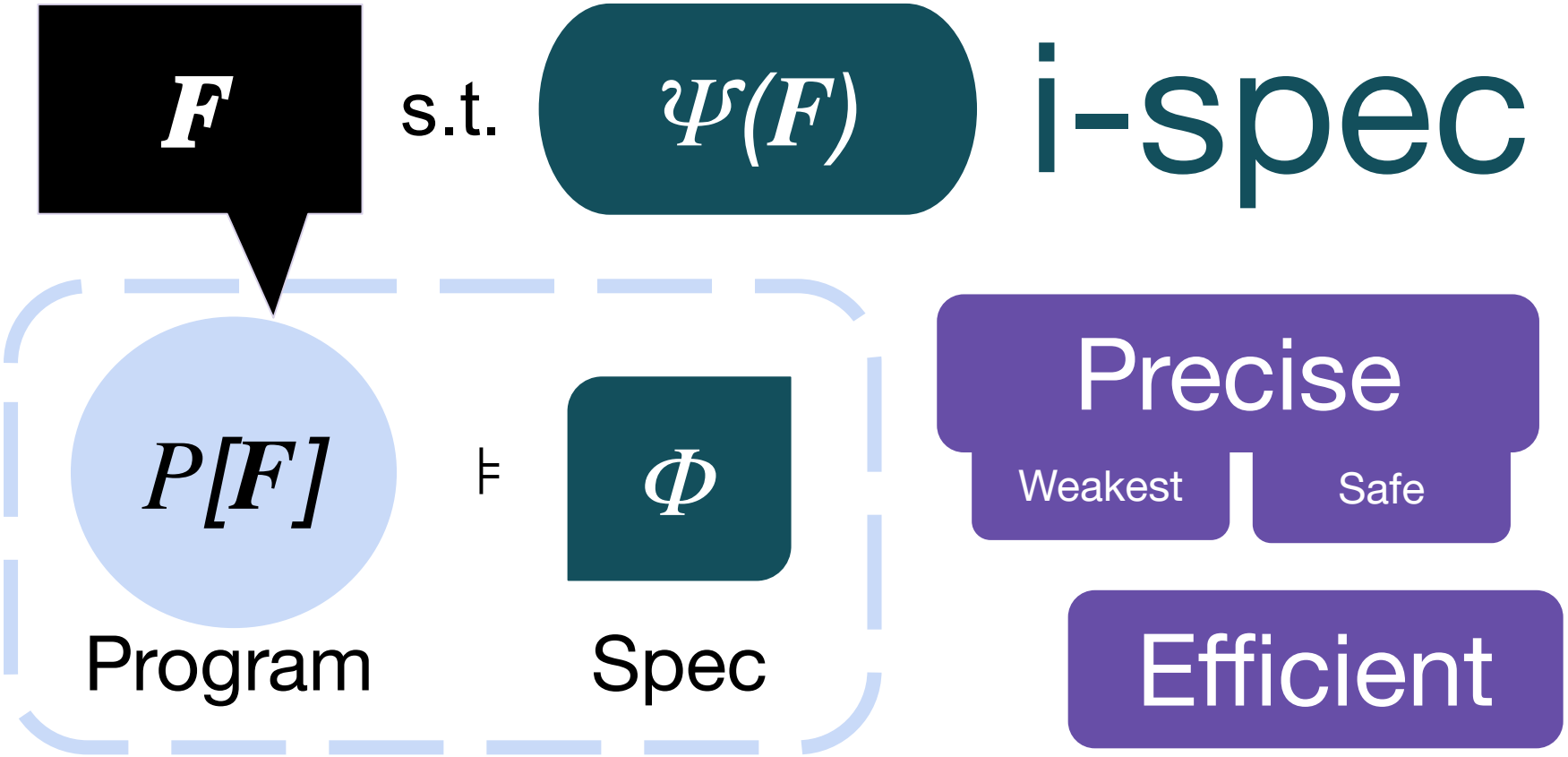
Not Safe!!!

Eliminates
no “good runs”
[VMCAI ‘13]

[POPL ‘16]

[SIGCOMM ‘20]


F
s.t.
Ψ(F)
i-spec
P[F]
⊨
Φ
Program
Spec
Safe
Weakest
Efficient


F
s.t.
Ψ(F)
i-spec
P[F]
⊨
Φ
Program
Spec
Precise
Weakest
Safe
Efficient

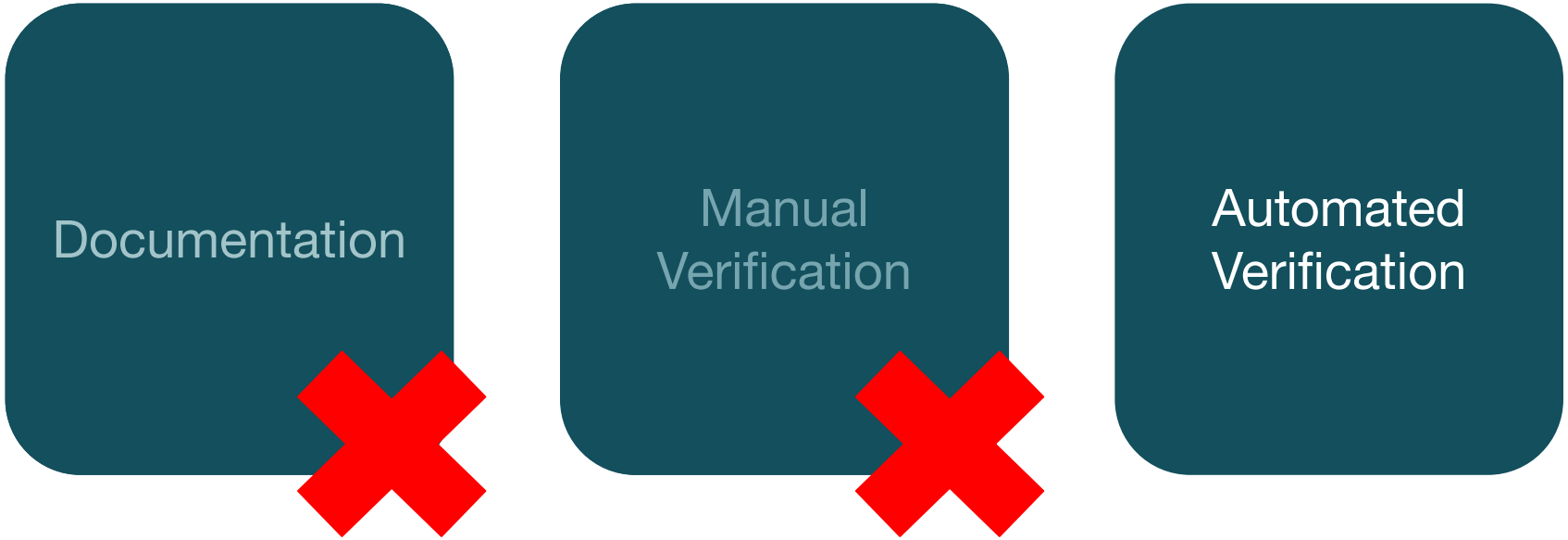
```
import unknown as foo,bar

def code(x):
    y = foo(x)
    z = bar(x)
    assert y|z ≠ 0x00
```

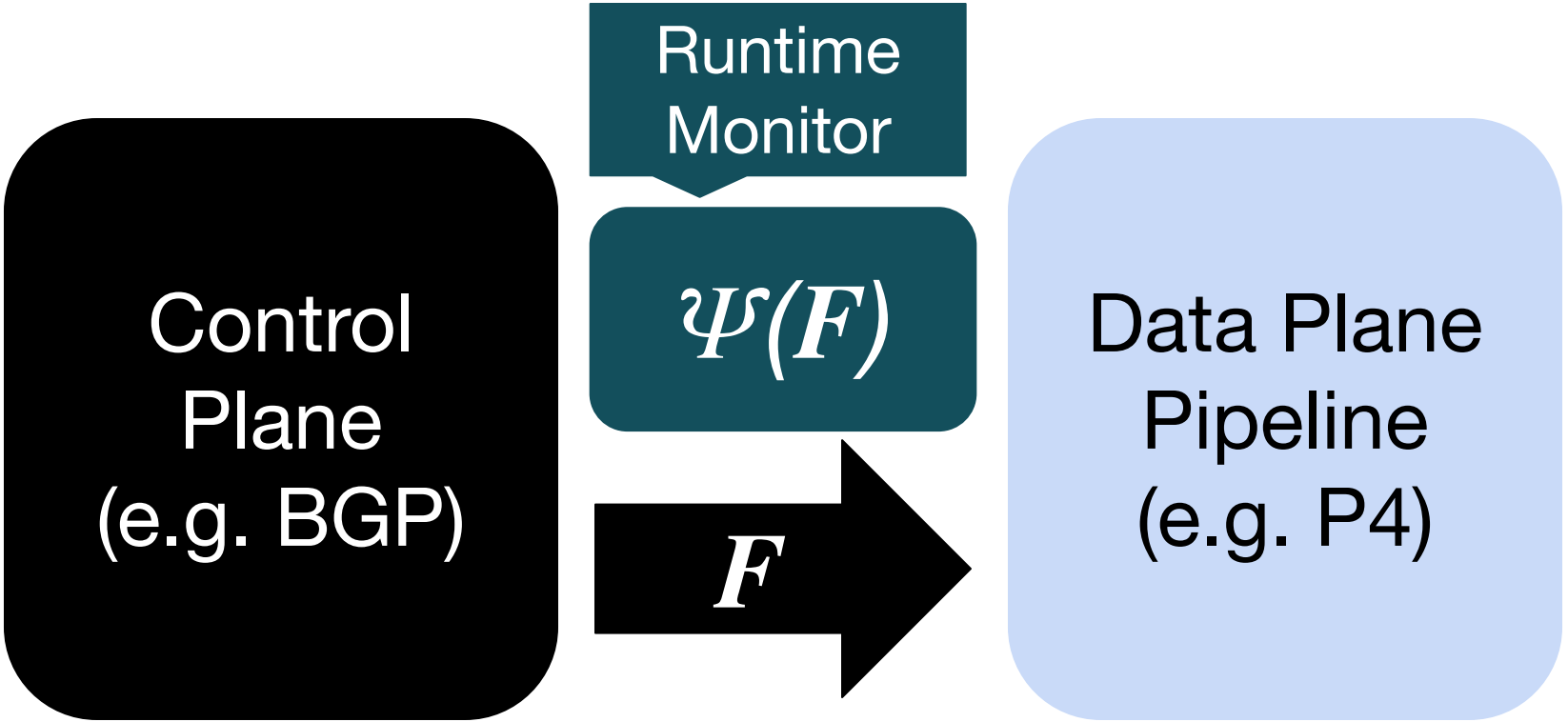
Goal:
Compute
precise i-specs

```
foo(x) | bar(x) ≠ 0x00
```

Safe | Weakest | Efficient

How to use computed i-specs?





Computer Network

Capisce

computes control interface specs
(ci-specs)

Precise

Safe

Weakest

Efficient

Step 1: Model Pipeline Programs

GCL(**F**)


```
c ::= assume φ
    | x := e
    | c ; c
    | c [] c
```

F : $2^w \to 2^l$

e ::= … | **F**(e)

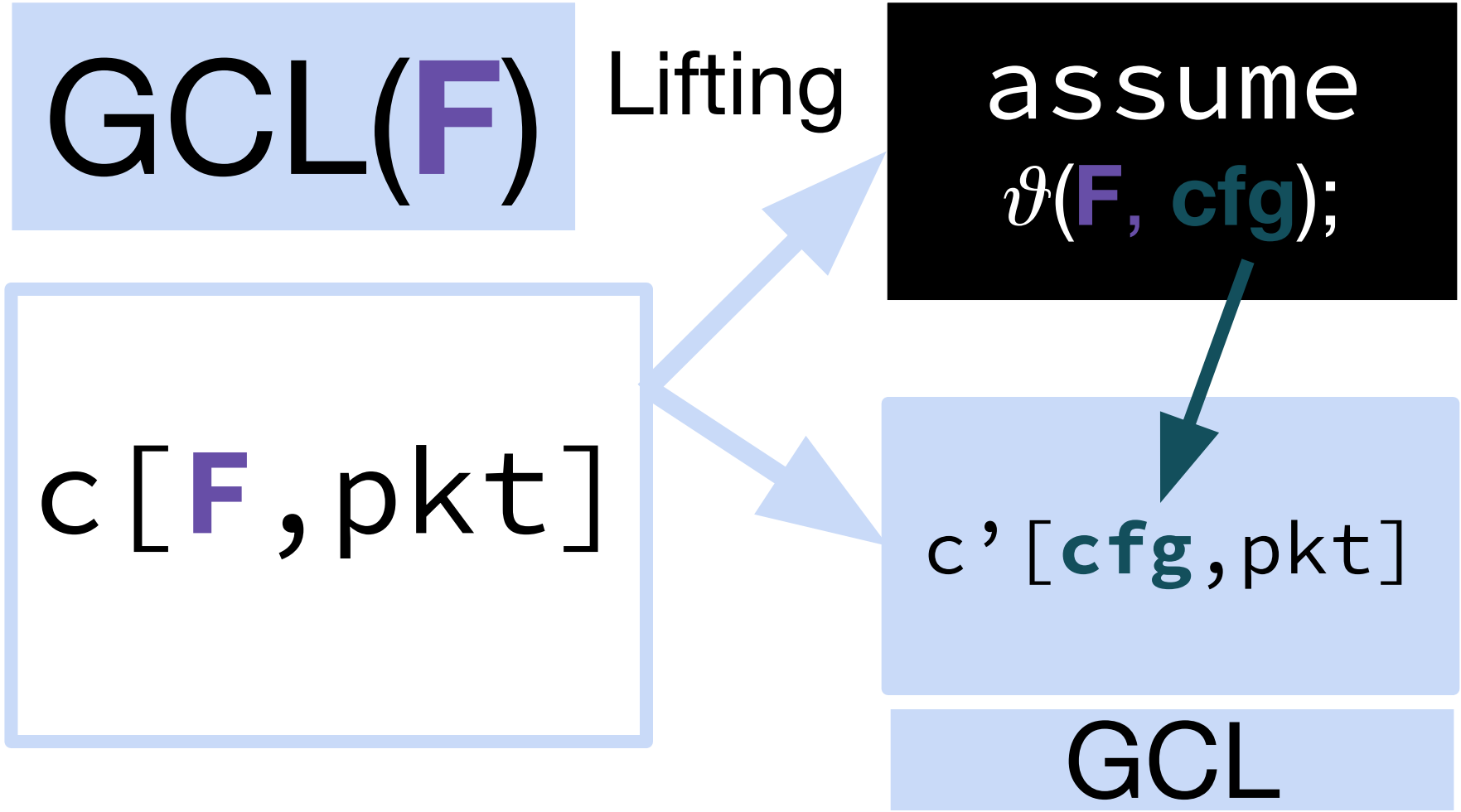
$\varphi \in$ QFBV

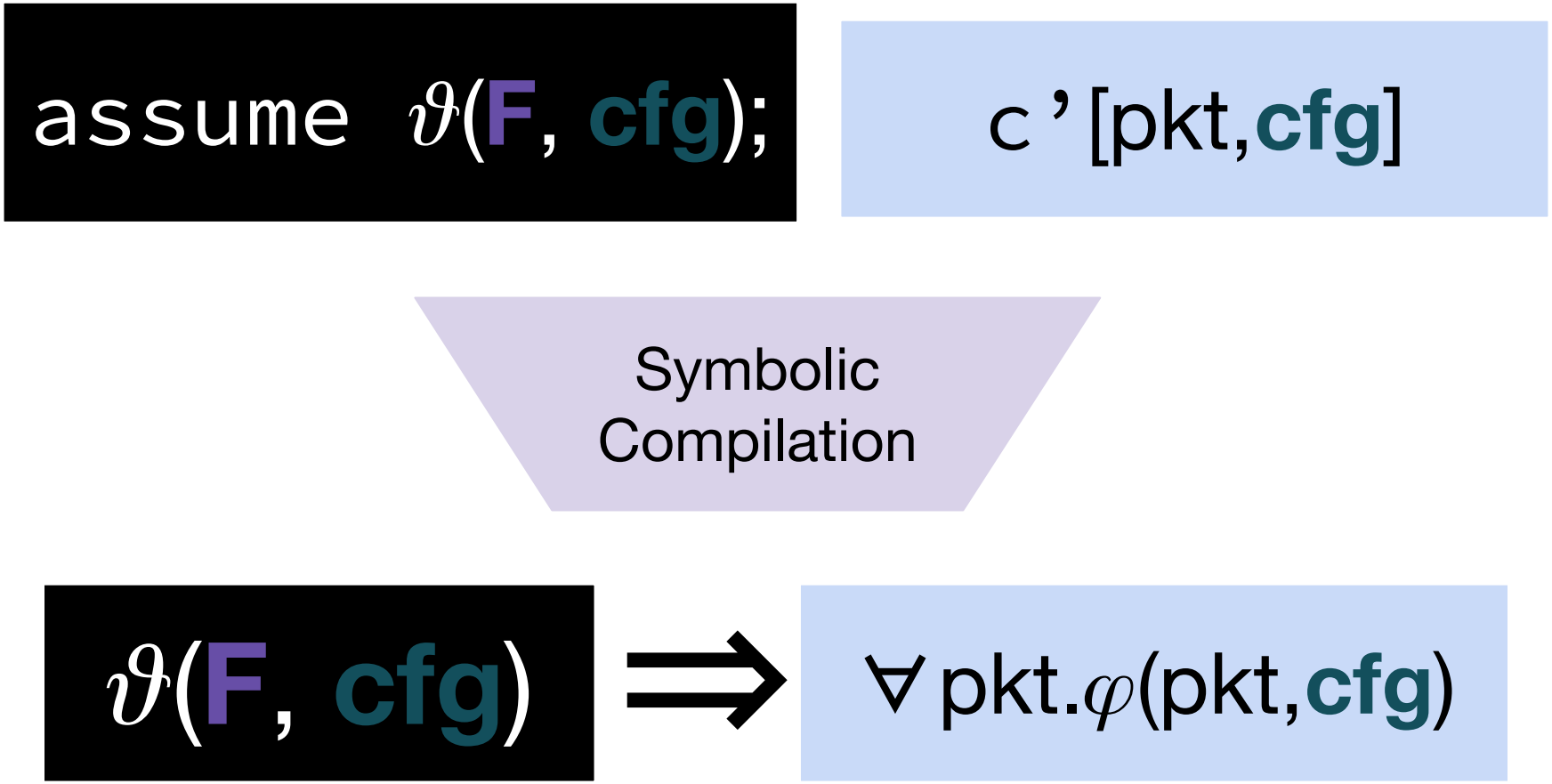
Control Flow in GCL(F)

$\texttt{if}\ \varphi\ c_1\ c_2$

$\texttt{assume}\ \varphi; c_1$
$[]$
$\texttt{assume}\ \neg\varphi; c_2$

Step 2: Symbolic Compilation


GCL(F)
Lifting
assume
ϑ(F, cfg);
c[F,pkt]
c'[cfg,pkt]
GCL


assume ϑ(F, cfg);
c ' [pkt,cfg]
Symbolic
Compilation
ϑ(F, cfg)
$\Rightarrow$
$\forall$ pkt.φ(pkt,cfg)

ci-spec

ϑ(F, cfg) $\Rightarrow$ $\forall$ pkt.φ(pkt,cfg)

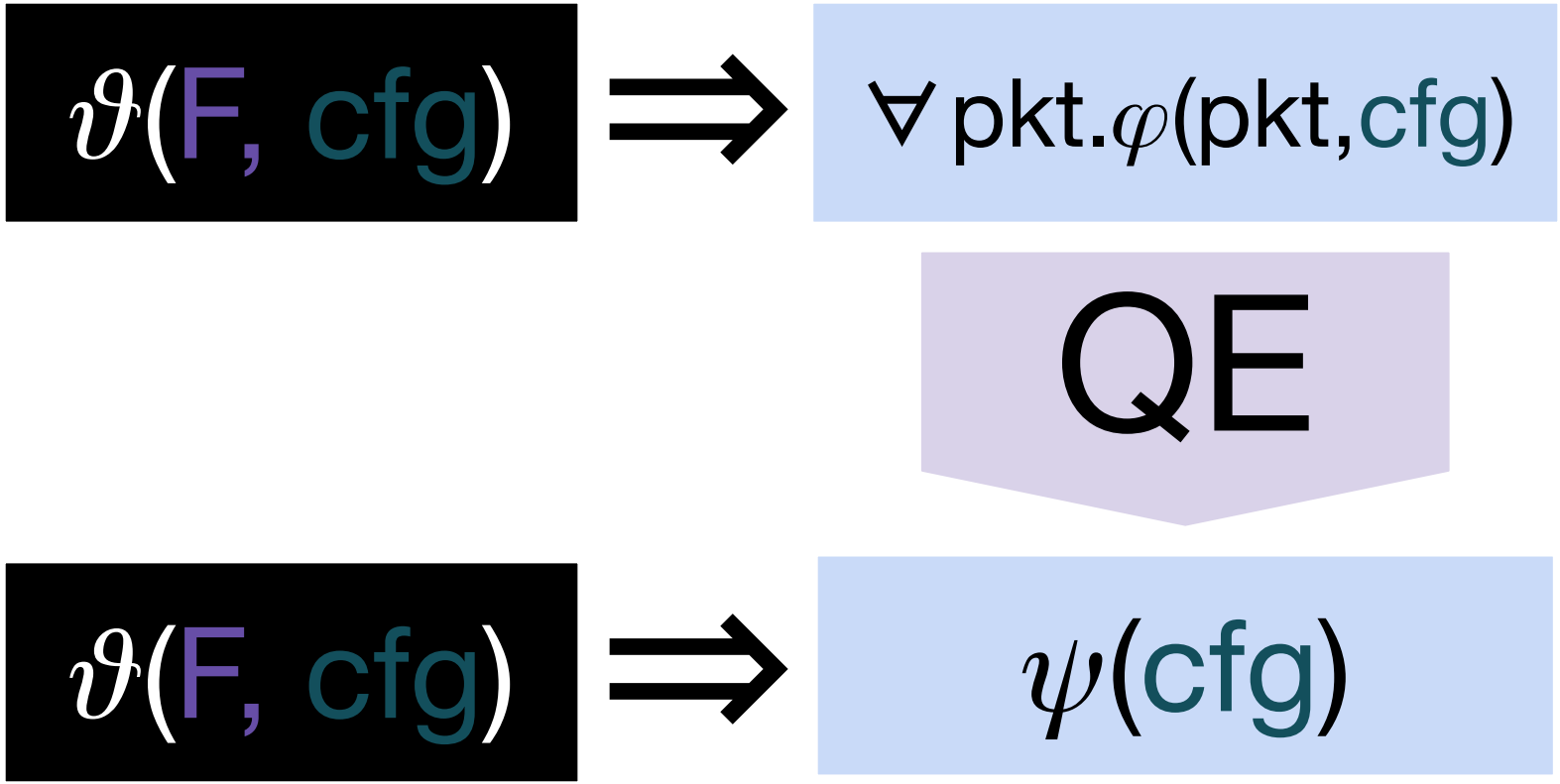
~1500 bits

pkt is big

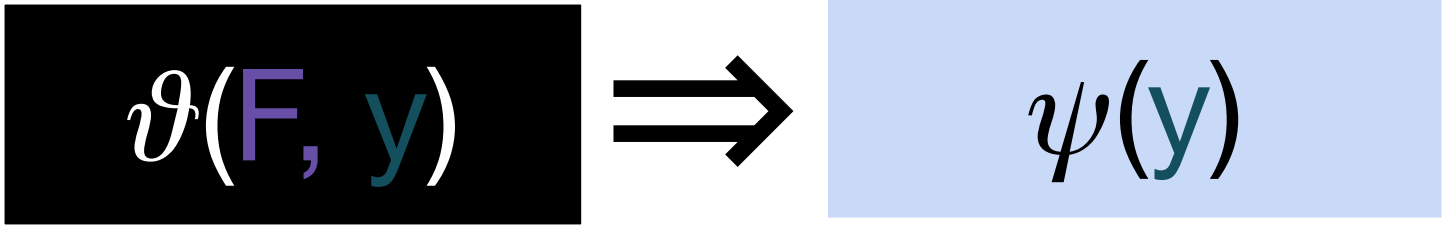
Precise

¬Efficient to monitor

Step 3: Quantifier Elimination


ϑ(F, cfg)
$\Rightarrow$
$\forall$ pkt.φ(pkt,cfg)
QE
ϑ(F, cfg)
$\Rightarrow$
ψ(cfg)

ci-spec



Theorem. *Precise*
=> *safe*
=> *weakest*

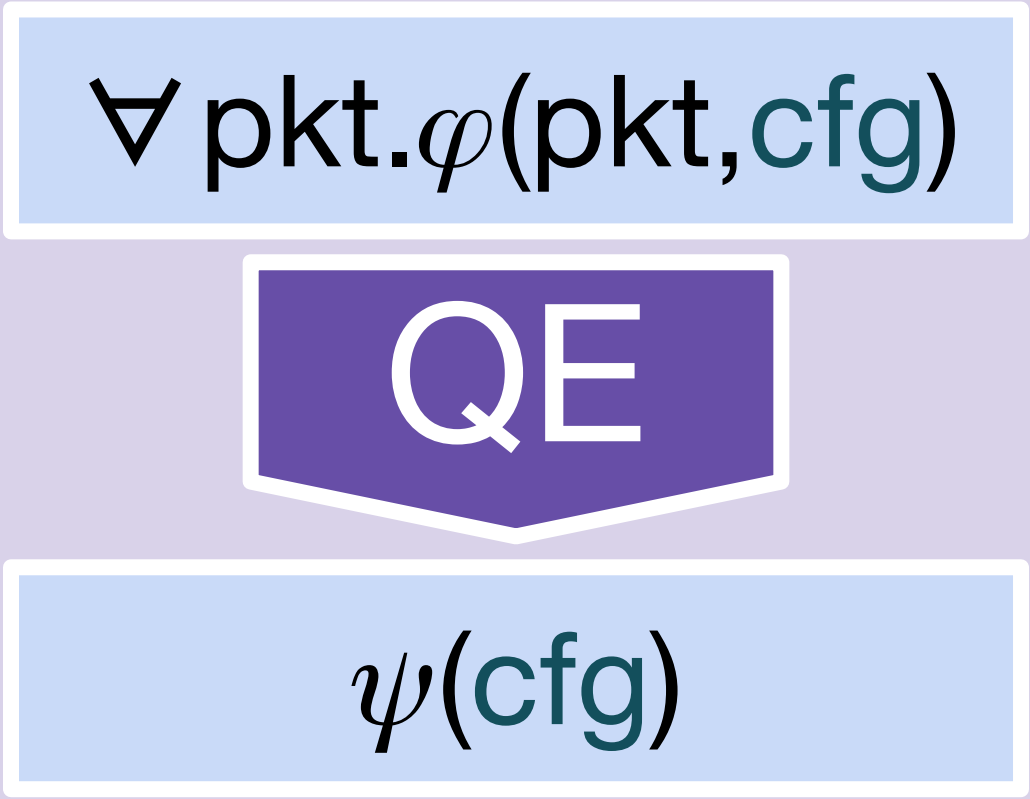
Theorem. *Terminates*

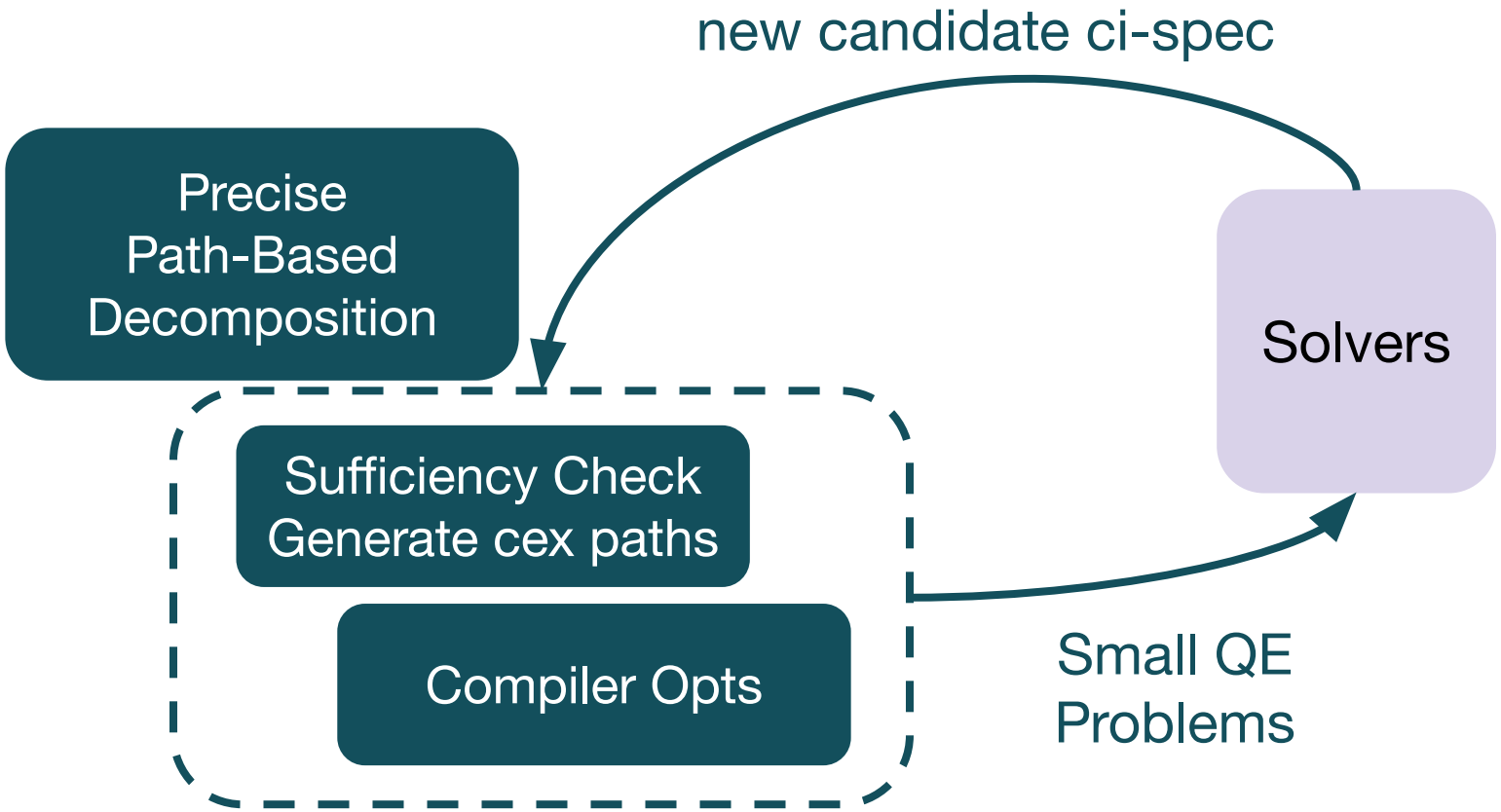
Theorem. *Efficiently monitorable*

Efficiently Control-Monitorable Sentences

… have polynomial expression complexity

Quantifier Elimination is *Intractable*




new candidate ci-spec
Precise
Path-Based
Decomposition
Solvers
Sufficiency Check
Generate cex paths
Compiler Opts
Small QE
Problems

Evaluation

Survey of Industrial and Academic P4 Programs

Ensure Invalid Data is not Read

Program	Program Paths	Result	Time (s)	Explored Paths	Spec AST Size	Explored Ratio
			Absurd Programs			
ts-switching	21	⊥	0.160	2	1	0.095
mc-nat	39	⊥	0.089	1	1	0.026
			Fixes to Absurd Programs			
ts-switching-fixed	21	⊤	0.030	0	1	0.0
mc-nat-fixed	39	⊤	0.027	0	1	0.0
			Trivial Programs			
resubmit	9	⊤	0.028	0	1	0.0
netpaxos-acceptor	0.116	⊤	30.0	0	1	0.0
ecmp	102	⊤	0.030	0	1	0.0
hula	3629	⊤	0.068	0	1	0.0
ndp-router	3843	⊤	2.9	0	1	0.0
			Nontrivial Programs			
arp	95	φ	5.0	0.016	349	0.17
heavy-hitter-2	267	φ	0.29	3	26	0.011
heavy-hitter-1	327	φ	0.60	7	90	0.021
flowlet	649	φ	1.8	9	127	0.014
simple_nat	66531	φ	5.2	54	1421	0.00081
07-multiprotocol	54459	φ	16	143	3138	0.0026
netchain	26726780	φ	2.9×10^3	264	11658	9.9×10^{-6}
linearroad	54477696		timeout			
fabric	133365047559893		timeout			
			Spec Smell Programs Fixes			
heavy-hitter-1-fixed	327	φ	0.63	7	107	0.021
linearroad-fixed	54477696	φ	5.9×10^4	3236	179885	5.9×10^{-5}
fabric-fixed	133365047559893	φ	1.2×10^3	653	41140	4.9×10^{-12}

Program Survey

Ensure Invalid Data is not Read

Program	Program Paths	Result	Time (s)	Explored Paths	Spec AST Size	Explored Ratio
Absurd Programs						
ts-switching	21	$\bot$	0.160	2	1	0.095
mc-nat	39	$\bot$	0.089	1	1	0.026
Fixes to Absurd Programs						
ts-switching-fixed	21	$\top$	0.030	0	1	0.0
mc-nat-fixed	39	$\top$	0.027	0	1	0.0
Trivial Programs						
resubmit	9	$\top$	0.028	0	1	0.0
netpaxos-acceptor	0.116	$\top$	30.0	0	1	0.0
ecmp	102	$\top$	0.030	0	1	0.0
hula	3629	$\top$	0.068	0	1	0.0
ndp-router	3843	$\top$	2.9	0	1	0.0
Nontrivial Programs						
arp	95	φ	5.0	0.016	349	0.17
heavy-hitter-2	267	φ	0.29	3	26	0.011
heavy-hitter-1	327	φ	0.60	7	90	0.021
flowlet	649	φ	1.8	9	127	0.014
simple_nat	66531	φ	5.2	54	1421	0.00081
07-multiprotocol	54459	φ	16	143	3138	0.0026
				264	11658	9.9×10^{-6}
linearroad	54477696		timeout			
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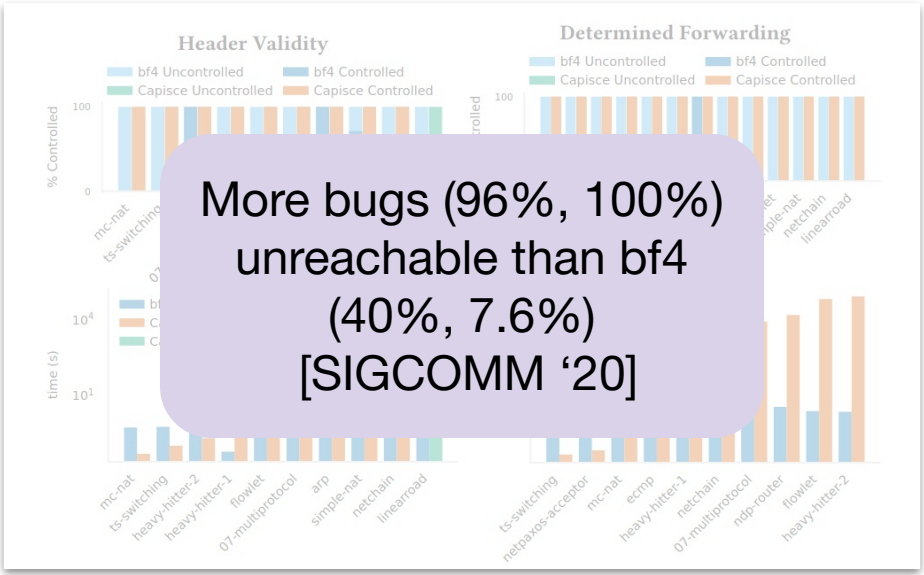
Spec Smell Program Fixes						
linearroad-fixed	54477696	φ	5.9×10^{4}	3236	179885	5.9×10^{-5}
fabric-fixed	133365047559893	φ	1.2×10^{3}	653	41140	4.9×10^{-12}

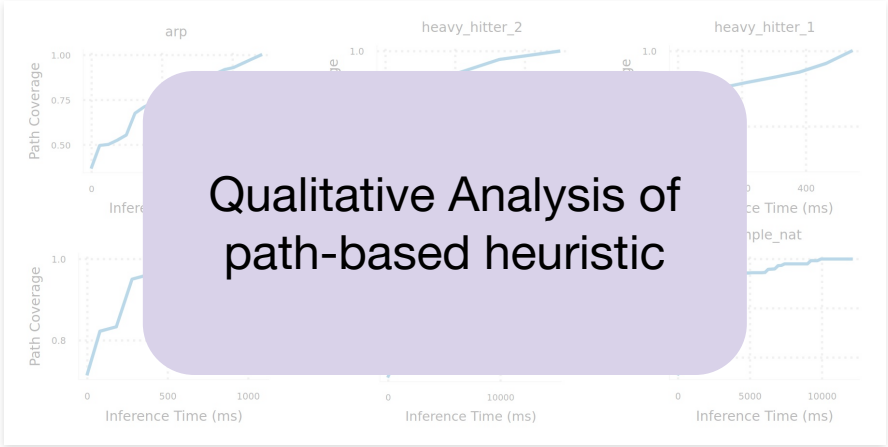
Program Survey

Defined Forwarding

Program	Program Paths	Result	Time (s)	Explored Paths	ci-spec Size	Explored Ratio
		Abs	rd Progra	s		
ecmp	102	⊥	0.320	4	1	0.039
fabric	133365047559893	⊥	7.3	5	1	3.7×10^{-14}
netchain	26726780	⊥	27	7	1	2.6×10^{-7}
		Triv	al Progra	s		
arp	95	⊤	0.027	0	1	0.0
linearroad	54477696	⊤	0.054	0	1	0.0
simple-nat	5548	⊤	0.034	0	1	0.0
		Nonti	vial Prog	ams		
resubmit	9	φ	0.016	2	17	0.22
ts-switching	21	φ	0.10	1	4	0.048
mc-nat	39	φ	0.27	3	21	0.077
netpaxos-acceptor	116	φ	0.12	1	4	0.0086
heavy-hitter-2	267	φ	88	15	233	0.056
heavy-hitter-1	327	φ	0.10	11	187	0.034
flowlet	649	φ	79	15	490	0.023
hula	3629	φ	0.39	1	9	0.00028
ndp-router	3843	φ	40	36	824	0.0094
07-multiprotocol	54459	φ	30	232	5034	0.0043
		Spec	mells & Fi	es		
ecmp-fixed	102	φ	0.28	3	34	0.029
mc-nat-fixed	27	⊤	0.029	0	1	0.0



More bugs (96%, 100%) unreachable than bf4 (40%, 7.6%) [SIGCOMM ‘20]



Qualitative Analysis of path-based heuristic

Capisce

$\Psi(F)$

p[F]

Φ